



Försättsblad till skriftlig tentamen vid Linköpings Universitet



Datum för tentamen	2015-01-09
Sal (1)	<u>TER3</u>
Tid	8-13
Kurskod	TDDD72
Provkod	TEN1
Kursnamn/benämning Provnamn/benämning	Logik En skriftlig tentamen
Institution	IDA
Antal uppgifter som ingår i tentamen	4
Jour/Kursansvarig Ange vem som besöker salen	Tommy Persson
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Besöker salen ca klockan	ja
Kursadministratör/kontaktperson (namn + tfnr + mailaddress)	Anna Grabska Eklund, ankn. 2362, anna.grabska.eklund@liu.se
Tillåtna hjälpmedel	1. You can use your own copies of slides as well as an English-Swedish dictionary. 2. Exercises are formulated in English, but answers can be given in English or Swedish.
Övrigt	
Antal exemplar i påsen	

EXAM: TDDD72 (LOGIC)

9 JANUARY 2015

RULES

1. You can use your own copies of slides as well as an English-Swedish dictionary.
2. Exercises are formulated in English, but answers can be given in English or Swedish.
3. You are not allowed to:
 - use any writing material other than indicated in point 1;
 - use calculators, mobile phones or any other electronic devices;
 - lend/borrow/exchange anything during the exam.
4. If an exercise has not been specified completely as you see it, state which (reasonable) assumptions you have made.
5. Begin each exercise on a new sheet of paper. Write only on one side of the paper. Write clearly and make sure to give adequate explanations for all your answers.
6. There are 4 exercises, each exercise gives maximum 10 points (40 points together). Grading is provided in the following table.

number of points (n)	Swedish grade	ETCS grade
$34 \leq n \leq 40$	5	A
$27 \leq n < 34$	4	B
$20 \leq n < 27$	3	C
$n < 20$	not passed	F (not passed)

EXERCISES

EXERCISE 1

1. Prove the following propositional formula

$$[(\neg R \vee P) \wedge (\neg(P \vee Q) \vee R) \rightarrow (\neg Q \vee P)]$$

- (a) (2 points) using tableaux;
 (b) (2 points) using Gentzen system (as provided in the book or during lectures – up to your choice).
2. Prove the following formula of predicate logic, where a, b are constants:

$$\left[\forall y Q(b, y) \wedge \forall x \forall y [Q(x, y) \rightarrow Q(s(x), s(y))] \right] \rightarrow \exists z [Q(b, z) \wedge Q(z, s(s(b)))].$$

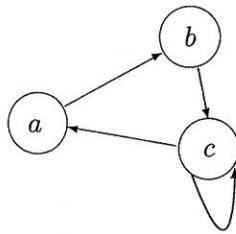
- (a) (3 points) using tableaux;
 (b) (3 points) using resolution.

EXERCISE 2

1. (4 points) Translate the following claims of L, T into a set of propositional formulas, where L, T and John are persons:
- claims of L :
 - “John rests on Saturdays.”
 - “John reads books and does not watch TV on Saturdays.”
 - claims of T :
 - “when John does not rest, he does not watch TV either.”
 - “on Saturdays John reads books or watches TV or cooks.”
2. (2 points) Assuming that L always lies and T always tells the truth, hypothesize what is John’s activity.
3. (4 points) Prove your claim formally using a proof system of your choice (tableaux, Gentzen system or resolution).

EXERCISE 3

Consider games, in which players move stones between places. A move from place p to place q is possible if there is an arrow between p and q . For example, in the game represented in the following figure,



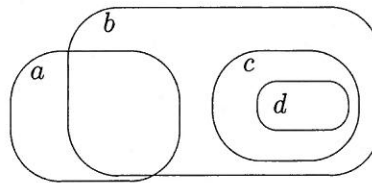
there are moves from a to b , from b to c , from c to a and from c to c and there are no moves from b to a , from c to b , from a to a and from b to b .

Different games are obtained by drawing different connections between places. A particular game can, for example, satisfy the following conditions:

- (a) "From every place there is a move."
 - (b) "Every two places p, q have the property that whenever there is a move from p to q then there is also a move from q to p ."
 - (c) "For every places p, q, r , if there are moves from p to q and from p to r then there is also a move from q to r ."
1. (3 points) Express in predicate logic properties (a), (b) and (c).
 2. (2 points) Check informally whether the conjunction of (a), (b) and (c) implies that "From every place there is a move to itself".
 3. (5 points) Verify your informal reasoning using resolution.

EXERCISE 4

1. (2 points) Design a Datalog database for storing information about regions on a map. Each region is characterized by its size (small, medium or large) and its characteristics (industrial, agricultural or nature). In addition, the database should contain information about *direct subregions*, denoted by \sqsubset . For example, in figure below we have regions a, b, c and d , and we have $c \sqsubset b, d \sqsubset c$ but it is not the case that $d \sqsubset b$ or $a \sqsubset b$.



2. (1 point) Express in predicate calculus the constraint:
 “every region has a unique (exactly one) size and a unique characteristics.”
3. (1 point) Provide an exemplary integrity constraint concerning the relation \sqsubset .
4. Formulate in logic queries selecting:
- (2 points) all large industrial direct subregions of region a ;
 - (4 points) all small (not necessarily direct) subregions of medium agricultural regions, where region a is a *subregion* of b if $a \sqsubset b$ or there is $n \geq 1$ and regions r_1, \dots, r_n such that $a \sqsubset r_1 \sqsubset r_2 \sqsubset \dots \sqsubset r_n \sqsubset b$.