

1467

2017-08-21

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TODDB68

TEN 9

Multiple choice form for answering question 1. Please put X:s in the appropriate cells:

	A	B	C	D	
1 a)				X	1
1 b)	X	X			1
1 c)		X			1
1 d)	X		X		0
1 e)	X			X	0
1 f)	X		X		0
1 g)	X		X		1
1 h)	X	X			1
1 i)			X	X	0
1 j)	X			X	1

6p

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2

```

a) void fx (int A) {
    wait (condX) // separate these
    // Body of fx
    *A = return value
    signal (condX)
}

```

(dereference)

This ensures that the functions are mutually exclusive. OK

```

b) wait (condX) {
    while (test_and_set (condX, false) == false) {
        while (condX == false)
    }
    condX = false
}

```

Semaphore implementation

```

int signal (condX) {
    if (condX == false) {
        condX = true
        return 0
    } else {
        return 1
    }
}

```

[Handwritten signature]

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2

b) The test and set function is atomic and in this case sets `convX` to false while returning its previous value. The while loop after it is just there for performance to limit I/O by limiting writes.

c) The current implementation has no race condition so no changes are needed. ~~X~~

1.5p

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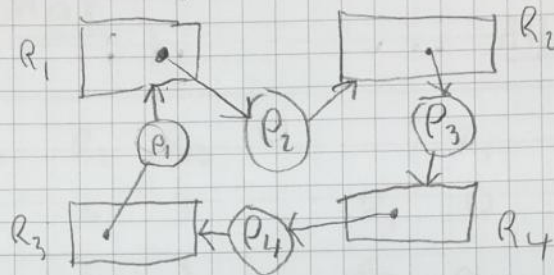
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Deadlocked

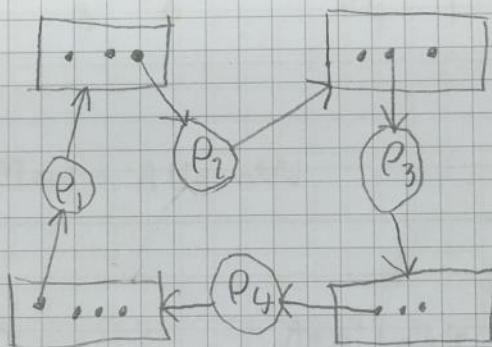
a)

(i)



ok

(ii)



No deadlock

ok, but not 4 instances.

2,5p

b)

3

b) Max

Allocated

	R ₁	R ₂	R ₃		R ₁	R ₂	R ₃
P ₁	0	6	3	P ₁	0	1	2
P ₂	0	5	3	P ₂	0	5	3
P ₃	1	0	6	P ₃	0	0	0
P ₄	1	3	6	P ₄	1	2	0

Need = Max - Allocated

Need

Available = [4, 0, 1]

	R ₁	R ₂	R ₃	
P ₁	0	5	1	≤ [4, 0, 1] == false
P ₂	0	0	0	≤ [4, 0, 1] == true Available := [4, 5, 4]
P ₃	1	0	6	≤ [4, 5, 4] == false
P ₄	0	1	6	≤ [4, 5, 4] == false

We need a new iteration over the need matrix.

Need

	R ₁	R ₂	R ₃	
P ₁	0	5	1	≤ [4, 5, 4] == true Available := [4, 6, 6]
P ₂				No need to do for P ₂ again
P ₃	1	0	6	≤ [4, 6, 6] == true Available := [4, 6, 6]
P ₄	0	1	6	≤ [4, 6, 6] == true Available := [5, 8, 6]

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b) Yes, this request is safe to grant and we get the safe sequence:

$$P_2 \rightarrow P_1 \rightarrow P_3 \rightarrow P_4$$

Ok. You need to check
request \leq available.

Also would be good if you explain which step you are doing.

3.5p

5p

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4 a) A process is an active entity executing its program. A user level thread does the same thing, i.e. executing code and following its instructions. The difference comes in how they are used, a process can have many user-level threads executing different parts of the programs code.

A process can exist without any user level threads, a user level thread must always belong to a process.

A kernel-level thread is used to map the user level threads and processes to a processor. The processor can only schedule kernel-level threads so mapping to them is essential. While user-level threads are thought to be parallel, if they map to the same kernel thread which is blocked, they both get blocked (both user-level threads).

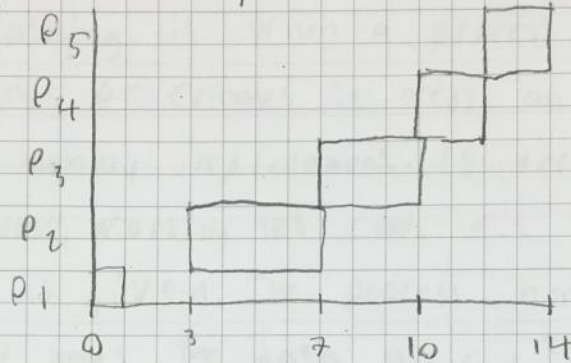
user-level

All threads belonging to the same process share many data structures where as two different processes share no data by default.

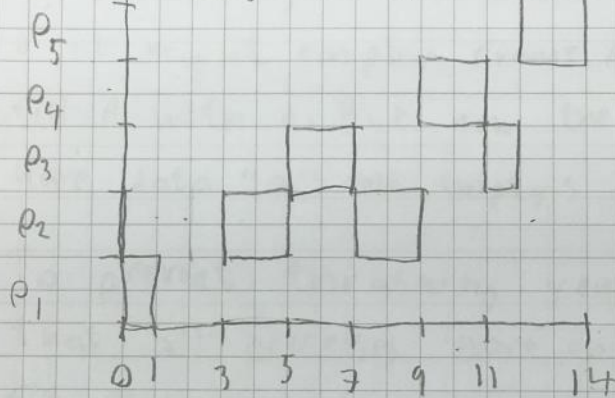
4

b)

(i) FIFO

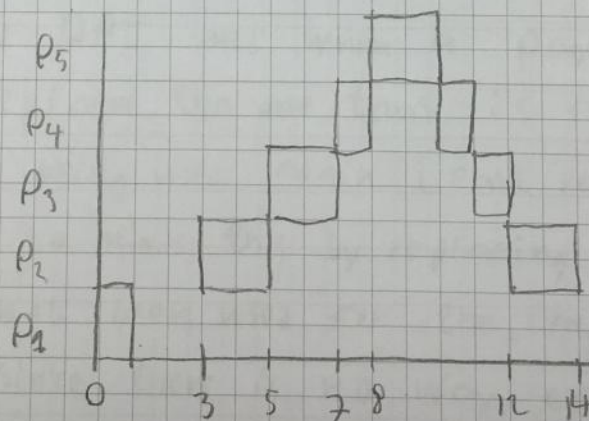


(ii) Round Robin



priority scheduling with preemption

(iii)



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a) Thrashing is when a process replaces the content of frames it needs, an endless loop of memory replacement. It occurs when the process' working set can't fit into the memory it is given. The process has 3 frames A, B, C and needs the data a, b, c, d. If A contains a, B contains b, C contains c. To fit the data d it might empty frame A of data "a" and fill A with d. But now the process doesn't have data "a" and empties another frame. And so on...

To prevent thrashing you can make sure that all processes have enough memory to fit their working set. (2)

b) The optimal page replacement algorithm is called OPT and when a page is to be replaced it replaces the one that is furthest away from being used again. Least recently used (LRU) tries to mimic this by replacing the page that has not been used for the longest time. It does this since there is no way of knowing which page will be used next unless we can see the future. (2)

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6 a) The reason for opening a file is to find it more easily. All open files are saved in a list so to find them we just go through this list. If this was not the case every operation called on a file would have to find it first, wasting time.

The open systemcall adds the file to the kernel's list of open files and returns a handle for the file.

②

b) Having a single copy of the file saves space and allows for instant updates if the file changes. If the file is shared through an acyclic graph system the system would need to continuously check for cycles with for example a garbage collector, which takes time and resources.

If the file is used by many users who often want write access to the file, having a single file would likely cause conflicts.

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b) Letting users have their own version of the file is the default implementation in Linux and windows, so this implementation is likely easier than the other.

1.5